QUALITY OF SERVICE STABILITY BASED MULTICAST ROUTING PROTOCOL FOR MANETS

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ABSTRACT


KEYWORDS: MANETS, QoS, TUNED CSA, DQSMRP

1. INTRODUCTION

Mobile Ad hoc Networks (MANETs) are foremost in class kind of wireless communication networks in which mobile nodes link on an on-demand basis. MANETs are having the ability to configure on their own and easing peer-level communications among mobile nodes independent by having no fixed infrastructure. When the route traffic occurs, constructing a MANET with each node’s to continuously retain the information remains a major challenge. Normally, MANETs are used for group communications where multicast protocols are efficiently compared to unicast protocols since they improve the efficiency of the wireless links in MANETs. The multicast protocols are also used as application demands for transmitting copies from either single or multiple sources to the receivers.

Multicasting reduces the communication costs by sending the single copy of the data to multiple recipients through which it minimizes link bandwidth, processing and transmission delay. Previous researchers have developed different algorithms to improvise the ability for selecting the routing path(s) and thereby, effectively met the desired Quality of Service (QoS) [1]. This infrastructure consequences in a highly dynamic topology causing a challenging task for QoS Routing.

When compared with well-established networks such as Wi-Fi, Global System for Mobile communication (GSM) and Code Division Multiple Access (CDMA), MANET was not reliable to meet QoS performances. The MANET performance is determined by the QoS through the basic parameters like bandwidth or delay or loss. QoS routing addresses both the issues of deciding a path from source to destination and fulfilling the QoS constraints like bandwidth, delay...
and loss. Route selection method is used to fulfill QoS requirements. The goal of QoS routing strategy is to maximize the network resources in order to reduce the constraints in MANET’s mobile nodes.

In real time applications, QoS routing protocol works based on multi-hop mobile networks [3, 4]. Particularly, QoS requirement depends on constraints such as link constraints, path constraints and tree constraints. End to End delay is usually affected by path constraint whereas, delay-jitter [5] by tree constraint. In order to meet all the constraints, robust techniques are required to solve the QoS requirements. Heuristic techniques are better to solve these constraints rather than deterministic methods. Researchers have studied various improvements based on heuristic methods [5, 6] and at the same time devised various solutions to solve QoS routing problems using heuristic methods [7].

Ad hoc mobile routing (AM Routing) was studied through which a specific unicast routing protocol was identified as more feasible than other unicast routing protocols (8). In QoS routing, network topology was constructed in which bandwidth link works in destination point [9]. A new protocol was proposed for ad hoc networks which are used to establish and maintain a shared mesh for each multicast group called Policy-based Unified Multi-radio, Architecture for Agile Mesh Networking (PUMA) [10]. A new single mixed metrics have been delineated for multi-constraint QoS routing [11]. Previous literature has found that Agent Based Multicast Routing Scheme (ABMRS) on MANET’s could be used as a set of static and mobile agents for QoS routing [12, 13]. Optimized link state routing protocol was launched and scrutinized in different research papers [14].

Besides that, several heuristic methods are examined to study the multi-constraint QoS routing. In [15], Simulated Annealing (SA) was introduced to solve the QoS routing problems. Reference [16] discussed Bees Life Algorithm (BLA) to solve the QoS multicast routing problem and constraints such as delay, allowed jitter, and requested bandwidth have been studied. Similarly, a Tabu search [17] based method concentrates on two important constraints in QoS such as bandwidth and end-to-end delay. Ant algorithm optimization was globally accepted for the effective and systematic handling of QoS routing [18][19] Harmony search (HS) algorithm, a unique algorithm was used to find the path of the MANET and it was discussed based on bandwidth-delay-constraints in multicast routing problem [20].

Hybrid methods have been proposed to address the delay of premature convergence [21]. A PSO-GA algorithm was evolved using a combination of both particle swarm optimization (PSO) method and genetic algorithm (GA) which provides an effective and efficient search for the solution. In [22], the new fuzzy genetic algorithm is also introduced for QoS multicast routing is also discussed.

Thus in this paper, we propose Dynamic Quality of Service Stability based multicast routing protocol (DQSMRP) for the multi-constraint QoS routing problem. Recently, Cuckoo Search (CS) is derived from reproduction strategies [23] which are widely used in heuristic based algorithms. [24, 25]. It is more exhaustive and provides quick search to obtain better solutions to routing in more reliable and robust manner.

2. PROBLEM DEFINITION, RESEARCH METHODOLOGY

The optimization problem is used to solve the QoS routing problem which is the primary objective to find a multicast tree constraints based on the cost function in order to reduce the practical constraints. Four different constraints are considered for formulating the current study
such as delay, packet loss, bandwidth and jitter. The following equation provides formulation data
for computation [5] QoS routing function is formulated as below

\[ \text{Min} C(H(x, s)) = \delta_1 C_d + \delta_2 C_b + \delta_3 C_d + \delta_4 C_p \] .........................................(1)

Reasons to merge both cost optimization and multi-constrained routing are as follows

- Formulations are based on multiple constraints and the multicast structure in order to spot
  an attainable path for each node having source and destination.
- To shrink the network resource utilization by users.

In progressively,

\[ C(H(x, s)) = \sum_{y \in H(x, s)} c(y) \] ........................................ (2)

The Delay Constraint is defined as the allowed delay limit that exists in any branch of the
network. When there is no link provided between any two nodes, a penalty occurs.

\[ \text{Delay}(R(x, y)) = \sum_{y \in R(x, y)} d_l(x) + \sum_{y \in R(x, y)} d_l(y) \] .........................(3)

The Bandwidth Constraint is the requirement of average bandwidth among all tree branches in the
network.

\[ \text{Bandwidth}(R(x, y)) = \min(bw(x), x \in R(x, y)) \] .........................(4)

Specifically \( \delta_1, \delta_2, \delta_3, \delta_4 \) are the bandwidth, delay, delay-jitter and packet-loss for the
corresponding penalty constants. The objective of the initiated solution is to solve the multicast
QoS routing problem as described by numerical experiments in the coming section.

2.1 TUNED CUCKOO SEARCH MULTICAST ROUTING PROTOCOL

A technique that produces simple implementation and quality solutions is the best solution
provider. This paper provides better results over the existing algorithms such as meta-heuristic
techniques. Before analyzing the proposed method, the cuckoo algorithm is discussed as follows.

2.1.1. AN OVERVIEW OF CUCKOO SEARCH ALGORITHM

By analogy of reproduction strategy, CS algorithm is a well known prominent heuristic search
algorithms which follow the principle of cuckoos (chosen nests of other birds for laying eggs)
[26, 27]. Generally, the host bird would not permit the cuckoo eggs to easily differentiate and
vice-versa. The important optimization procedure is as follows:

- **Initial Solution**: Cuckoo egg symbolizes a set of solutions and its value of
dimensions was extracted from various nests randomly;
- **Next Generation**: The best eggs with common solutions only able to proceed to the next
  generation.
- **Acceptance Rule**: New egg will succeed in the nest if eggs are unknown. Important step
  of CS are explained as follows:
• Step 1: Randomly set “Pop” host nests, a group of the population. (e.g., \( P = \sum_{i=1}^{P_{op}} P_i \))
• Step 2: Evaluate the fitness of all nests \( P = \sum_{i=1}^{P_{op}} f(P_i) \)
• Step 3: Do Iterations <Max. Iterations.
• Step 4: Produce a cuckoo egg \( P^j \) by applying a Lévy flight from a random nest
• Step 5: Compute the fitness of \( f(P_j) \)
• Step 6: Check for existence: choose again a random nest \( J \)
• Step 7: if \( f(p^j_i) < f(P_i) \) then
• Step 8: \( P_i = p^j_i \)
• Step 9: end if
• Step 10: Remove a fraction of the worst nests at the rate of \( R_d \)
• Step 11: Randomly, Building new nests at new locations using Levy flights replacements
• Step 12: End While

The benefit of this heuristics is its simplicity. The single parameter of CS algorithm should be considered enough when compared with PSO and Ant colony algorithms.

### 2.1.2 Node Stability

For efficient better packet delivery services in forwarding group, the mobile node is to be very stable. The way of providing stability is that moving node around its current position. Node stability metric have been used in previous work of authors is given in [28]. It is studied that stable path nodes have been found from packets from the source to the multicast group. Self-stability and neighbor nodes stability are two metrics to represent node stability.

Following steps to find stability of a node:

- In MANET’s nodes, it establishes self-stability, i.e., node movement is relative to its previous position.
- MANET’s node gives the node stability factor from self-stability and neighbor stability. Actually, stability is based on the distance between the movements of nodes positioning within transmission range. Stable node is created when movement within transmission range.

### 2.1.3 Neighbor Node Stability

It can be defined as how the node is being connected to its neighbour to obtain self-stability. Whenever nodes are in transmission range, it can exchange messages with each other. Each node consists of connectivity information, the signal stability of one-hop neighbors, and also maintains a neighbor list. The degree of a node \( n \) is denoted as number of links (or nodes) which connected to it and denoted as ND.

\[
N_s(t) = \alpha \times \frac{1}{ND} \sum_{i=1}^{ND} S_{ij}(t) + (1 - \alpha) \times N_s(t - 1)
\]

\[\text{.........} (5)\]

Where, \( \alpha \) is the weighting factor (lies between 0 and 1) and is distributed between 0.6 and 0.7 since they yield better results in simulation.
2.1.4 **LINK STABILITY**

Link stability between the nodes defines quality and existence of the connection. It is reckoned based on two parameters:

- Received Signal Strength and
- A lifetime of the Link.

The algorithm represents a pseudocode for updating link stability status between the nodes.

The different parameters used in the algorithm are as follows:

- **Lifetime** – Duration of continuous connectivity between the nodes
- **Lifetime threshold** – Indicates the maximum limit of link lifetime that decides link stability
- **Link stability status** – It is a Boolean variable that defines link stability between the nodes.
- **Recent** – Indicates most recent response received for a Hello packet from a neighbor
- **P** – Represents Total Number of Hello packets
- **Received signal strength** – Represents the strength of the signal received from a neighbor
- **Signal threshold** – Is acceptable signal strength to be received from neighbors

### 2.1.5 ALGORITHM: LINK STABILITY STATUS BETWEEN THE NODES

- **Step 1**: \( P = \text{No of Hello Packets}; \)
- **Step 2**: \( \text{Lifetime} = 0; \)
- **Step 3**: Link stability status = 0;
- **Step 4**: Recent = 0;
- **Step 5**: Lifetime threshold = \( P \times \text{Hello Packet Interval}; \)
- **Step 6**: While \( P > 0 \) do
  - **Step 7**: If received signal strength \( \geq \) signal threshold then
    - **Step 8**: Lifetime = lifetime + 1;
    - **Step 9**: Recent = 1;
    - **Step 10**: \( P = P - 1; \)
    - **Step 11**: else
    - **Step 12**: Recent = 0;
    - **Step 13**: \( P = P - 1; \)
  - **Step 14**: end if
- **Step 15**: end while
- **Step 16**: lifetime sec = lifetime \( \times \) Hello Packet Interval;
- **Step 17**: if (lifetime sec \( > \) lifetime threshold) and (Recent) then
  - **Step 18**: link stability status = 1;
  - **Step 19**: else
  - **Step 20**: link stability status = 0;
  - **Step 21**: end if
2.1.6 Delay Estimation

For delay estimation, an arbitrary node that contributes to traffic forwarding using the M/M/1 queuing system is modeled. This queue represents a single queuing station with a single server [29]. The authors assume that the contributing nodes are served by a single server with first come first serve queuing policy. Packets arrive according to a Poisson distribution rate (λ), and the probability distribution of the service rate is exponential, denoted by μ. The maximum size of the queue in every node is represented by K.

To satisfy delay requirements in multimedia real-time applications, packets must be received by multicast receivers which satisfy the application delay constraints. When a packet is to be sent either by a source node or forwarding group of nodes, it experiences three types of delays: queuing, contention and transmission delay. The total delay considered over a link between two nodes is given by

\[ d_{Total} = d_q + d_c + d_t \]  

The queuing delay denoted by \( d_q \) is the delay between the time the packet is assigned to a queue and the time it starts transmission. The packet waits and gives other packets in transmission queue are transmitted at this particular time. This is the amount of time a packet is spent in the interfacing queue (T).

2.1.7 Bandwidth Estimation

Bandwidth is one of the key principles for assessing Quality of service (QoS). The authors considered their previous work presented in [30], to calculate the available bandwidth based on their channel status of the radio link to calculate the idle and busy periods of the shared wireless media. Bandwidth is measured by the node activities and its surrounding neighbors using the channel utility.

In IEEE 802.11 MANETs, once it is achieved to gain the channel access, a node can be able to transmitting data packets. Hence, a node senses the channel and estimates bandwidth by using the idle and busy times in a predefined interval. This is expressed by the following equation.

\[ BW = \frac{T_{idle}}{T_{interval}} \times \lambda \]  

2.1.8 Route Discovery Process using Tuned Cuckoo Process

Multicast stable QoS path creation involves two phases: a request and a reply phase. Request phase invokes route discovery process to find routes having a group of receivers using stable and QoS intermediate nodes. It helps to reply phase for updating of RIC and also performs confirmation process. It acts as intermediate nodes that help to create a multicast mesh to the group of receivers.

RR packets are used in source node to find the route to its group of receivers. The sequences of operations that occur are as follows:

- Source node prepares an RR packet with an application of bandwidth and delay requirements.
Selective transmission of RR packet to neighbors once satisfies stability criteria, i.e., SFBN greater than SFTH, and bandwidth requirement, i.e., estimated bandwidth greater than twice the application requirements.

Once received, node discards RR packet (both sequence number and source address).

If RR packet is not a duplicate, check if the availability of Routing Information Cache (RIC); if available, RP packet create and start reply propagation to the source.

If RR packet is a duplicate, then discard it and stop transmission of RR packet.

If not duplicate and no route available in RIC, transmit the RR packet by updating its fields (route record, Stability Factor Between Nodes (SFBN ) record, bandwidth record, delay record, time to live, and nexthop address) to its neighbors as in step 2.

Perform steps 3 to 6 (Till destination reached).

If the receiver is not reached within certain hops, send RE packet to the source node.

3. RESULTS AND DISCUSSION

3.1 SIMULATIONS AND PERFORMANCE EVALUATION

In this section, the performance of proposed protocol with DQSMRP and Fully Distributed Multicast Routing Protocol (FDMRP) [31] is compared, through a considerable set of simulations. These etiquettes have been taken for comparison because both are mesh based. These protocols are collated in terms of packet delivery ratio, control overhead and average end-to-end delay. Simulations considering the values of the production parameters are taken for several iterations, and these are computing the mean. The values lying within 95% of the confidence interval of the mean are used for computing the mean value, which is shown in the graphs in result analysis section. The various network scenarios have been simulated. Simulation environment consists of four models: Network, Channel, Mobility and Traffic. An ad hoc network is created at an area of l × b square meters as network model; it is having N number of mobile nodes spaced randomly. Limited bandwidths in coverage area around each node have shared to its
neighbors. It is assumed that the operating range of transmitted power and communication range are constant.

### 3.1.1 PERFORMANCE PARAMETERS

Following metrics have been used to scrutinize the performance:
- **Packet Delivery Ratio (PDR)**: The ratio of a number of average data packets received at the multicast receivers to the number of data packets sent by the source.
- **Packet Over Head**: The ratio of control packets sent to the network to the total number of average data packets delivered to the receivers.
- **Average End-to-End Delay**: The average delay experienced by the successfully delivered packets in reaching their receiver.

<table>
<thead>
<tr>
<th>No. of Nodes</th>
<th>50,100,150 and 200.</th>
</tr>
</thead>
<tbody>
<tr>
<td>Routing Protocol</td>
<td>DQSMRP</td>
</tr>
<tr>
<td>Area Size</td>
<td>1000 X 1000</td>
</tr>
<tr>
<td>Mac</td>
<td>802.11</td>
</tr>
<tr>
<td>Radio Range</td>
<td>250m</td>
</tr>
<tr>
<td>Simulation Time</td>
<td>10 sec</td>
</tr>
<tr>
<td>Traffic Source</td>
<td>CBR</td>
</tr>
<tr>
<td>Packet Size</td>
<td>512</td>
</tr>
<tr>
<td>Receiving Power</td>
<td>0.395</td>
</tr>
<tr>
<td>Sending power</td>
<td>0.660</td>
</tr>
<tr>
<td>Idle Power</td>
<td>0.035</td>
</tr>
<tr>
<td>Initial Energy</td>
<td>10.0 J</td>
</tr>
<tr>
<td>Node Speed</td>
<td>5, 10, 10 and, 20 ms</td>
</tr>
<tr>
<td>No of destinations</td>
<td>2, 4, 6 and 8</td>
</tr>
</tbody>
</table>

| Table 1: Simulation scenario |

### 3.2 RESULT ANALYSIS

The simulation results of DQSMRP and FDMRP present the comparison results by varying number of multiple destinations and node speed. According to the simulation results in Fig 2 to Fig 7, the performance of TCSMRP is more reliable than FDMRP. The execution of DQSMRP packet delivery ratio is gathered more than 14.2% compare to FDMRP. Moreover, the rendering of DQSMRP starts to end delay is less than FDMRP. The performance rate is almost 18% less and the delay ratio is varied with respective of node speed and number of destinations. According to the overhead performance scenario, the overhead rate is 22% less than FDMRP.
Figure 2: PDR vs Number of Destinations

Figure 3: PDR vs Node Speed
**Figure 4:** Overhead vs Node Speed

**Figure 5:** Overhead vs Number of destinations
4. CONCLUSIONS

This paper clarifies about a hybrid computational intellectual algorithm i.e., Tuned Cuckoo Search Algorithm (TCSA), that includes the dominant characteristic of two different heuristic techniques have been used to solve QoS multicast routing. The QoS routing in a MANET is a
intricate to determine a prime routing. The general conclusion from presented simulation experiments reveals that proposed routing protocol performs better than FDMRP in terms of packet delivery ratio, packet overhead and average delay from start to end as a function of varying number of receivers, sources and node speeds. In future works, the author’s aims to study further by comparing our Tuned Cuckoo Search Multicast Routing Protocol with some more QoS based routing protocols in MANETs. Numerical experiments have shown that the propounded hybrid TCSA algorithm provides advanced search multicast trees, particularly in both swift convergence and robustness, when compared to other algorithms. This proposed algorithm shall be the escort for real-time networks with suitable sophistication based on when and where it utilized.

REFERENCES