MOBILITY LOAD BALANCING BASED ADAPTIVE 
HANDOVER IN DOWNLINK LTE SELF-ORGANIZING 
NETWORKS

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ABSTRACT

This article investigates mobility load balancing (MLB) algorithm implementation through network simulator (ns-3) in long term evolution (LTE) systems employing orthogonal frequency division multiple access (OFDMA) for downlink (DL) data transmission. MLB is introduced by the third generation partnership project (3GPP) as a key target of LTE self-organizing networks (SONs) [1]. Our contribution is twofold. First, we implemented elementary procedures (EPs) related to load management (LM) function of the X2-application protocol (X2AP) as specified in TS 136.423 [2]. We particularly focused on EPs 'Resource Status Reporting Initiation Procedure' and 'Resource Status Reporting Procedure'. Second, we implemented a MLB based adaptive handover (HO) algorithm enabling to configure adaptively HO hysteresis threshold for each neighbouring cell, of an overloaded cell, according to its current load information. Numerical results show how, through suitable simulation scenarios, MLB enables enhancing network performance in terms of overall throughput, packet loss ratio (PLR) and fairness without incurring HO overhead.

KEYWORDS

LTE, load management, X2AP, elementary procedure, mobility load balancing

1. INTRODUCTION

In recent years, wireless cellular networks’ operators have experienced a spectacular growth in mobile data traffic. With the surging traffic demand, wireless cellular networks are becoming significantly more complex, causing higher operational costs. In order to reduce these costs while optimizing network efficiency and service quality, the SON concept had been introduced in LTE systems. SON concept aims to reduce manual operations by integrating self-optimizing, self-healing and self-configuration features. MLB is one of the most important functionalities that belong to self-optimizing SON techniques.

In wireless cellular networks, loads in different cells are frequently unequal. This leads to critical situation where hot spot cells (i.e. overloaded cells) suffer from a high call blocking rate (CBR) and/or call dropping rate (CDR). In contrast, a large part of resources in low-loaded cells remains in idle state causing resource wastes. Therefore, load imbalance between neighbouring cells seriously deteriorates network performance. In this regard MLB solutions are deployed to circumvent this problem and to enhance fairness, availability, scalability, user-experience and network utilization.

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In literature several load balancing methods devoted to cellular networks have been proposed. Authors in [3] formulate a multi-objective optimization problem using utility function jointly optimizing load balancing (LB) index and network average load for users with quality of service (QoS) requirements. They conclude that better values of LB index and network average load can be achieved when compared with other conventional methods. In [4, 5], authors investigate conflicting and stand-alone SON functions. They particularly focused on two SON functions namely MLB and handover optimization (HOO). In this direction also and in order to optimize the same parameter values having conflicting goals at the same network element (NE), 3GPP proposes an additional entity, usually called coordinator [21]. Alternatively, to reach the above purpose, authors in [4] propose a unified self-management mechanism for MLB and HOO based on fuzzy logic and reinforcement learning techniques. The objective behind is to reduce SON coordination entity’s complexity. Authors in [5] propose an optimization technique of both MLB and HOO functions based on a weighted co-satisfaction factor (CSF). Another practical solution for MLB is proposed in [6] and is based on HO margins’ auto-tuning of each base station according to its load and the load of its neighbouring cells. Authors highlight through simulation analysis MLB impact on both call admission rate and user throughput. Z. Huang et al. propose in [7] a multi-traffic load balance algorithm (MTLB) to promote an efficient load balance. Yang et al. propose in [8] an auto-configuration technique of handover hysteresis threshold according to the load information of the current overloaded cell and its neighbours. The proposed technique triggers the MLB procedure according to the load condition of a cell and significantly reduces CBR, HO dropping/blocking rate and ping-pong HO rate. However, the paper neglects the impact of the proposed mechanism on some key performance metrics such as the global throughput in the network and the LB index. The proposed model in the above paper is implemented using a formal description technique based on a specification and a description language (SDL). Thereby algorithm efficiency in [8] needs to be improved for more realistic conditions and in this regard a more realistic implementation with respect to the 3GPP specifications could be established. In this paper we propose a MLB algorithm implementation based on discrete-event network simulator ns-3 [17]. To that end two implementation steps are followed. In a first step two LM EPs of the of X2AP protocol namely ‘Resource Status Reporting Initiation Procedure’ and ‘Resource Status Reporting Procedure’ are implemented as specified in TS 136.423 [2]. We exploit next these EPs to implement in a second step the MLB algorithm proposed in [8] based on adjusting adaptively HO hysteresis threshold according to cell load information. We specifically focus on A3 HO triggering event based on measurements of the reference signal received power (RSRP) [9]. Numerical results show how MLB enhances network performance without incurring significant overhead. We particularly show for DL how performance metrics such as global network throughput, PLR, number of successful HO and LB fairness index evolve for different user equipment (UE) densities.

The rest of this paper is organized as follows. Section II presents X2AP HO and LM signalling functions in 3GPP. We succinctly recall in section III some definitions related to MLB based adaptive handover for LTE Systems. In section IV we provide a simulation model description of the LTE system adopted in this paper. We propose in section V the X2AP EPs extension implemented in ns-3 before detailing in section VI, based on X2AP functionalities, a MLB algorithm implementation for LTE systems in DL. Finally, simulation results are presented in section VII for different KPIs metrics and concluding remarks are provided in section VIII.
2. HO AND MLB SIGNALLING FUNCTIONS OF X2AP IN 3GPP SPECIFICATIONS

In LTE architecture, the core network (CN) includes mobility management entity (MME), serving gateway (SGW) and packet data network gateway (PDN-GW), whereas E-UTRAN includes eNodeBs (eNBs) [10]. The X2 interface is defined between two neighbouring eNBs at two planes (Figure 1) [11]:

- **X2-UP (User-plane):** is the user-plane interface between two neighbouring eNBs. X2-UP protocol tunnels end-user packets between eNBs. The transport network layer is built on IP transport, and GTP-U is used on top of the UDP to carry the user-plane PDUs.

- **X2-CP (Control-plane):** is the control plane interface between two neighbouring eNBs. X2-CP has stream control transmission protocol (SCTP) as transport layer protocol and X2AP as application layer signalling protocol.

X2AP supports radio network layer signalling functions of the control plane between eNBs. 3GPP defines X2AP as a collection of EPs [2]. A X2AP EP is a unit of interaction between two eNBs. Since X2AP is responsible of signalling between two neighbouring eNBs, it ensures two major SON functions namely HO signalling functions (i.e. mobility management (MM), mobility robustness optimization (MRO), mobility parameters management (MPM) and load information exchanges function (i.e. LM)). The HO enables one eNB to hand over an user equipment (UE) to another eNB and requires the transfer of useful information to maintain the LTE radio access network (RAN) services at the new eNB. The LM function allows to exchange traffic load and its related signalling between neighbouring eNBs. In what follows we give more details about HO and LM functions.

2.1. HO Functions

In LTE the HO procedure is based on UE’s measurements; A serving eNB sends a connection reconfiguration message through the radio resource control (RRC) layer to inform an UE of its related measurement configuration. If the UE detects any event related to measurement configuration, a measurement event is triggered and the serving eNB is informed. The decision of triggering a HO is based on the detection of a special event. 3GPP specifies several reporting
events to trigger HO [9] depending on the measurements' values. The serving eNB should indicate to UEs the event to use in the connection reconfiguration message. In order to trigger a HO within LTE several kind of events may be adopted.

Our choice in this paper is directed toward A3 event, which might be triggered if a neighbouring cell measurement becomes offset better than the serving cell. The entering condition to be satisfied for A3 event can be simplified as follows:

\[ Mn > Ms + Hys + \text{Off} + \text{Freq} + S\text{Off} \]  

where \( Ms \) is the UE measurement corresponding to the serving cell and \( Mn \) is the UE measurement corresponding to the neighbouring cells. As part of the reporting configuration, the eNB can tell the UEs to measure either their cells RSRP, or their reference signal received quality (RSRQ) [12]. \( Hys \) is a hysteresis parameter for measurement reporting. If the UE sends a measurement report to its serving cell, then \( Hys \) prevents any more reports until the signal level changes by \( 2 \cdot Hys \) (Figure 2). Similarly, \( \text{Off} \) is a hysteresis parameter for HOs. If the measurement report triggers HO, then \( \text{Off} \) prevents the UE from moving back to the original cell until the signal level changes by \( 2 \cdot \text{Off} \). \( \text{Freq} \) (resp. \( S\text{Off} \)) is the difference between frequency specific offsets (resp. cell specific-offsets) of the serving and neighbouring cells [9]. The time to trigger (TTT) parameter is the time during which the condition of triggering A3 event needs to be maintained in order to trigger a measurement report (Figure 2).

If a measurement report is sent to the serving cell with a triggered A3 event, this latter may trigger a HO procedure by starting a MM function with the neighbouring cell. In what follows we describe the MM function of the X2AP protocol.

### 2.1.1. MM Function

The MM function covers four EPs [2]:

- Handover Preparation EP: during this procedure the serving cell initiates the HO with the HANDOVER REQUEST message and includes necessary information for neighbouring cell preparation to an incoming HO. The neighbouring cell responds back to the serving cell with a HO REQUEST ACKNOWLEDGE message (successful operation) or a HANDOVER PREPARATION FAILURE message (unsuccessful operation);
- SN Status Transfer EP: along this procedure the SN Status Transfer is sent from the serving cell to its neighbours to transmit uplink (UL) and DL packet data convergence protocol sequence number (PDCP SN) and hyper frame number (HFN) receiver status and transmitter status from the source eNB to the target eNB during an X2 handover. This is achieved for each evolved radio access bearer (E-RAB) for which PDCP SN and HFN status preservation applies [2];

- UE Context Release EP: By sending an UE CONTEXT RELEASE message, the target eNB informs the source eNB of HO success and a release of resources by the eNB source is triggered. The target eNB sends this message after receiving the PATH SWITCH ACKNOWLEDGE message from the MME. Upon reception of the UE CONTEXT RELEASE message, the source eNB can release data and C-plane related resources associated to the UE context. Therefore, any ongoing data forwarding may continue;

- Handover Cancel EP: is used by a source eNB to cancel an ongoing HO preparation or an already prepared HO.

2.1.2. MPM Function

MPM function covers one EP: the mobility settings change (MSC) EP. This latter enables an eNB to negotiate the HO trigger settings with a peer eNB controlling neighbouring cells [2].

2.1.3. MRO Function

MRO function detects and corrects automatically errors in the mobility configuration. It covers two EPs:

- Radio link failure (RLF) indication EP [2]: The purpose of this EP is to transfer information regarding RRC re-establishment attempts or received RLF reports, between eNBs;

- Handover report EP: This EP enables to transfer mobility related information between eNBs.

2.2. LM Function

In order to detect a load imbalance between eNBs in the network, it is necessary to compare cells’ load information and exchange them between neighbouring eNBs [13].

Three separate EPs are used to exchange load information via the X2 interface:

- Load Indication EP: The load indication procedure is used over the X2 interface for load and interference management information exchange. An eNB initiates this EP by sending LOAD INFORMATION message to a neighbouring eNB;

- Resource Status Reporting Initiation EP and Resource Status Reporting EP (Figure 3): These EPs are used to initiate and report load measurements result between eNBs. A source eNB sends a RESOURCE STATUS REQUEST message to a target eNB. If the latter is able to provide requested resource status information, it responds with the
RESOURCE STATUS RESPONSE message and then Resource Status Reporting EP could be initiated. The target eNB will report measurement results in RESOURCE STATUS UPDATE message.

3. MLB BASED ADAPTIVE HO IN DL LTE SELF-ORGANIZING NETWORKS

To improve the overall system capacity 3GPP defines a self-optimizing function namely MLB. MLB enables cells suffering from congestion to transfer extra-load to other neighbouring cells willing to cooperate toward their spare resources [14]. The MLB is governed by load reporting mechanisms activated between eNBs (over the X2 interface) to exchange information about load level and available capacity. One solution to introduce MLB in a LTE system is to steer extra-traffic toward neighbouring cells by optimizing dynamically the cell HO parameters, such as hysteresis, upon detection of an imbalance [13]. In fact, SON techniques allow to automatically adjust mobility configuration based on several factors such as received load information.

In order to describe MLB principle let’s consider event A3 (Figure 2) described above and assume that it is triggered between an overloaded cell and its neighbour at time $T1$. The decreasing of the $H_{ys}$ value fosters the triggering of A3 event ($T2 < T1$) and consequently facilitates the HO triggering from the overloaded cell to its neighbouring cell (Figure 4). By adjusting $H_{ys}$ value, the hot spot cell forces UEs to hand over to its neighbouring cell. Then MLB based adaptive HO may be defined as an explicit triggering of a forced HO for load balancing purposes. To apply this mechanism, each cell should be aware of its own load conditions as well as the load of its neighbours.

Figure 3. Resource Status Reporting EP
This enables respectively to detect an overload situation and to select the best neighbouring cell to hand over UEs. For that purpose, eNBs periodically monitor their own load conditions and exchange load information over the X2 interface (Figure 3). The RESOURCE STATUS UPDATE message includes a special attribute called composite available capacity (CAC) [2]. CAC represents the overall available resource level that can be offered for MLB purposes in either DL (CACD) or UL (CACU). CAC can be represented as a data structure including two information elements (IE) [2]:

- Cell Capacity Class Value IE: classifies the cell capacity with regards to the other cells. The Cell Capacity Class Value IE indicates resources that may be reserved for data traffic;

- Capacity Value IE: Indicates the ratio between the amount of available resources and the amount of total E-UTRAN resources. The Capacity Value IE is measured and reported so that suitable E-UTRAN resources of the existing services are reserved. The Capacity Value IE can be weighted according to the ratio of cell capacity class values.

In this paper we turn our interest only to CACD which is not supported by the eNB model in given in current version of ns-3 (i.e. ns-3.24) [17]. In this regard we will define and develop the calculation of CACD. In the next section we will describe the system model and explain basic set notation we used for the formulation of the treated problem.

4. System Model

We consider in this paper a scenario of a multi-cell LTE network (Figure 5) where each eNB consists of three cells. Let $C$ denotes the set of all cells in the LTE network. $C_i$ represents the cell $i$ ($i \in [1..C]$).
Figure 5. Network model

$C_i$ is used to indicate the serving cell. $C_{i,j}$ denotes a neighbouring cell $C_j$ of a cell $C_i$ ($j \in [1..C]$). $C_{i,S_j}$ indicates the set of all neighbours $S_j$ of $C_i$. The amount of resource blocks (RBs) quantifying time-frequency resources of $C_i$ is denoted $R_i$. (We assume that all cells have the same amount of RBs). $U$ is the set of all UEs scheduled in the network (we assume that $U$ is constant along the simulation duration), and $U_i(t)$ is the set of UEs scheduled in $C_i$ at time $t$ and $U_{i,k}(t)$ indicates an UE $k$ scheduled in $C_i$ at time $t$. $R_{i,k}(t)$ is the amount of consumed RBs by $U_{i,k}(t)$ at time $t$. $HYS_{i,j}$ is the hysteresis value of HO from $C_i$ to its neighbouring cell $C_{i,j}$. The load $\rho_i(t)$ of $C_i$ at time $t$ can be calculated as follows:

$$\rho_i(t) = \frac{\sum_{k \in U_{i,k}(t)} R_{i,k}(t)}{R_i}$$  \hspace{1cm} (2)

Then CACD of $C_i$ at time $t$ (i.e. $CACD_i(t)$) could be defined as follows [15]:

$$CACD_i(t) = 100 \cdot (1 - \rho_i)$$  \hspace{1cm} (3)

To measure network performance, we use the following metrics:

- network average load in DL $\zeta$: $\rho(t)$ is the network load at time $t$:

$$\rho(t) = \frac{\sum_{i \in C} \sum_{t} \rho_i(t)}{C}$$  \hspace{1cm} (4)

Let $T$ denotes the simulation duration and $p$ the epoch duration. Then $\zeta$ can be calculated as follows:

$$\zeta = \frac{\sum_{i \in C} \sum_{t} \rho_i(t)}{C \cdot \left(\frac{T}{p}\right)}$$  \hspace{1cm} (5)
- Jain’s fairness index [16]:
  
  \[ F(t) = \left( \frac{\sum_{i\in C} \rho_i(t)}{C \cdot \sum_{i\in C} \rho_i^2(t)} \right)^2 \in [0..1] \]  

  then the Jain’s fairness index \( F \) can be calculated as follows:

  \[ F = \left( \frac{\sum_{i\in C} \sum_{j\in T} \rho_i(t)}{C \cdot \sum_{i\in C} \left( \sum_{j\in T} \rho_i(t) \right)^2} \right)^2 \in [0..1] \]  

- Packet loss ratio (PLR) in DL: is defined as the difference between the number of transmitted packets and the number of received packets.

- Number of successful HO: Since ns-3 in its current version (i.e. ns-3.24) [17] doesn’t support the HO Failure procedure, the number of successful HO is indeed the number of triggering (or the triggering frequency) of a HO procedure. The following section discusses the proposed X2AP EPs extension proposed in this paper for ns-3.

5. THE PROPOSED X2AP EPs EXTENSION

Ns-3 [17] is an open-source discrete event network simulator, developed using C++ programming language. Ns-3 supports the simulation of 3GPP LTE cellular networks. Baldo et al. propose in [18] a simulation model for LTE HO scenarios using ns-3. In order to integrate SON functionalities, the proposed model implements X2 interface supporting RESOURCE STATUS UPDATE message. Such implementation offers only a prototype of this message and details regarding associated variables are not specified. In this paper we implement new added methods illustrated in Figure 6. This is achieved through two implementation steps. The first one concerns RESOURCE STATUS REQUEST and RESOURCE STATUS RESPONSE messages (i.e. Resource Status Reporting Initiation EP), which are necessary for the initiation of a RESOURCE STATUS UPDATE message (RESOURCE STATUS FAILURE message is not considered in this paper). Whereas the second focuses on RESOURCE STATUS UPDATE message (i.e. Resource Status Reporting EP). The added methods are detailed as follows:

- LteEnbRrc::DoSendResourceStatusRequest(): An overloaded cell executes this function to interrogate neighbouring cells about their load conditions.
- LteEnbRrc::DoRecvResourceStatusRequest(): This function is executed when a cell receives a RESOURCE STATUS REQUEST message.
- LteEnbRrc::DoSendResourceStatusResponse(): A cell that receives a RESOURCE STATUS REQUEST message will execute this method to send an acknowledgement to its neighbour.
- LteEnbRrc::DoRecvResourceStatusResponse(): This function is executed when a cell receives a RESOURCE STATUS RESPONSE message.
**6. MLB PROPOSED ALGORITHM**

In our proposed MLB algorithm we assume $C_i$ is overloaded if Cond. 1 is satisfied (Table 1). Indeed if $CACD_i$ is less than a given threshold $Th_{PreLB}$, then $C_i$ is considered as a hot spot (i.e. overloaded) cell and it enables the MLB algorithm. Now if Cond. 2 is satisfied, this means that $C_i$ is low-loaded enough to disable the MLB algorithm.

<table>
<thead>
<tr>
<th>Condition</th>
<th>Action</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Cond. 1:</strong> $CACD_i \leq Th_{PreLB}$</td>
<td>$C_i$ falls into the heavy load condition and the MLB algorithm is enabled.</td>
</tr>
<tr>
<td><strong>Cond. 2:</strong> $CACD_i \geq Th_{PostLB}$</td>
<td>$C_i$ falls into the low load condition and the MLB algorithm is disabled.</td>
</tr>
</tbody>
</table>

The proposed MLB algorithm is based on the adaptive modification of actual hysteresis value, between a serving hot spot cell $C_s$ and its neighbouring cells $C_{s,Sj}$, according to the neighbouring
cells’ load. Enabling the MLB algorithm means that $C_s$ is allowed to send RESOURCE STATUS REQUEST messages to $C_{s,j}$ (Figure 7). $C_s$ will update $HYS_{s,j}$ ($j \in S_j$) value only if $C_{s,j}$ answers with a RESOURCE STATUS RESPONSE message (Figure 3). If a status error occurs, then $C_{s,j}$ sends a RESOURCE STATUS FAILURE message and $C_s$ will not update the $HYS_{s,j}$ value (Figure 7).

![Figure 7. Sequence diagram of the proposed MLB algorithm](image)

The new hysteresis value between $C_s$ and $C_{s,j}$ (i.e. $Newhys_{s,j}$) is adjusted as follows:

$$Newhys_{s,j} = \alpha_{s,j} \cdot hys_{s,j}, \quad j \in S_j \quad (8)$$

Where $\alpha_{s,j}$ is a weighting factor in $[0..1]$ calculated as follows (Figure 8):

$$\alpha_{s,j} = \begin{cases} 1 & \text{if } CACD_j < ThAvail_{LB} \\ \frac{1}{ThAvail_{LB} - CACD_j} & \text{if } ThAvail_{LB} \leq CACD_j < ThPost_{LB} \\ 0 & \text{if } CACD_j \geq ThPost_{LB} \end{cases} \quad (9)$$

Figure 8 represents the $\alpha_{s,j}$ evolution with respect to $CACD_j$. 

99
7. PERFORMANCE EVALUATION

7.1. Simulation Setup

Table 2 summarizes the main parameters of the simulation scenario where a multi-cell LTE network composed of 7 eNBs (21 cells) in hexagonal layout (Figure 5) is considered. Distance between eNBs is assumed equals to 500 meters. In order to create an overloaded situation, we initially attached all UEs to one cell ($C_1$ of the central eNB). The mobility pattern considered is the Random Waypoint model [19, 20]. In such model nodes move independently to a randomly chosen destination with a randomly selected velocity. This allows to emulate a real life situation in a reasonable way. The number of UEs in the simulation scenarios can be changed by choosing a proper UEs density $\lambda$. We conducted a simulation campaign with different number of UEs ($U \in \{25, 50, 75, 100, 150\}$) in order to investigate load impact on performance of the proposed MLB algorithm.

Table 2. Simulation parameters

<table>
<thead>
<tr>
<th>Parameter</th>
<th>Value</th>
</tr>
</thead>
<tbody>
<tr>
<td>Cellular layout</td>
<td>7 eNBs (21 cells) in hexagonal layout</td>
</tr>
<tr>
<td>Inter-eNBs distance</td>
<td>500 m</td>
</tr>
<tr>
<td>Cell Tx power</td>
<td>46 dbm</td>
</tr>
<tr>
<td>Path loss model</td>
<td>$L = 128.1 + 37.6 \log(R)$</td>
</tr>
<tr>
<td>Channel fading</td>
<td>Typical urban</td>
</tr>
<tr>
<td>Carrier frequency</td>
<td>2 Ghz</td>
</tr>
<tr>
<td>System bandwidth</td>
<td>5 Mhz (25 RBs)</td>
</tr>
<tr>
<td>Traffic</td>
<td>Only control messages, no data traffic</td>
</tr>
<tr>
<td>Error model</td>
<td>None</td>
</tr>
<tr>
<td>UE distribution</td>
<td>$U \in {50, 75, 100, 125, 150}$ attached all to $C_1$ of the central eNB</td>
</tr>
<tr>
<td>UE movement pattern</td>
<td>30 dbm</td>
</tr>
<tr>
<td>UE measurement reporting interval</td>
<td>50 ms</td>
</tr>
<tr>
<td>Simulation duration</td>
<td>50 s</td>
</tr>
<tr>
<td>Time To Trigger (TTT)</td>
<td>256 ms</td>
</tr>
<tr>
<td>Hysteresis default value</td>
<td>3 db</td>
</tr>
<tr>
<td>Hysteresis margin with MLB</td>
<td>[0..3db]</td>
</tr>
</tbody>
</table>
7.2. Simulation Results

We investigate in this section numerical results related to MLB algorithm implementation. Figure 9 exhibits the LTE network global throughput in downlink. As observed, the MLB impact on network performance depends closely on UEs density. We distinguish two kinds of LTE network global throughput behaviour with respect to UEs density:

- The first one for relatively very low UEs density (U=25) and relatively very high UEs density (U=150) where MLB algorithm has low impact on the network performance. Indeed for U=25, the heavy load condition (Cond. 1) is very likely to be not verified (i.e. $CACD_i > ThPre_{LB}$, $\forall i \in C$) and the MLB algorithm will very probably not be enabled during the simulation. This explains why the network global throughput for such UEs density is almost the same with and without MLB algorithm activation (Figure 9).

- Whereas for U=150, almost all cells are very likely overloaded, then even if a cell $C_i$ falls under Cond. 1, $HYS_{i,j}$ ($j \in S_j$) will remain unchanged (i.e. $CACD_j \leq ThAvail \Rightarrow (a_{i,j} = 1)$) with and without MLB algorithm activation.

- The second one for $U \in \{50,75,100\}$ where MLB algorithm has significant impact on the network performance since it enhances the network global throughput (Figure 9). Actually in such case it is very likely that some network cells fall under the heavy load condition (Cond. 1). Also, a heavy loaded cell $C_i$ may decrease hysteresis values, $HYS_{i,j}$ ($j \in S_j$), of neighbouring cells willing to cooperate (i.e. $CACD_j > ThAvail_{LB} \Rightarrow 0 \leq \alpha < 1$) to enable MLB.

To sum up, the effect of the MLB algorithm and its efficiency on global network throughput are mainly tributary of UEs density and network cells load disparities.

![Figure 9. Global network throughput vs UEs density](image-url)
Figure 10 highlights a similar behaviour of Jain’s fairness index to the global network throughput (Figure 9) versus UEs density. In fact, the Jain’s fairness index values are improved when activating the MLB algorithm and this improvement is noticeable except for very low ($U=25$) and very high ($U=150$) network UEs densities.

Figure 10. Jain’s fairness index vs UEs density

Figure 11 represents the PLR evolution with respect to UEs density with and without activation of the MLB algorithm and shows that PLR is almost the same in both cases except when the total number of UEs in the network is around 100. Three kinds of PLR behaviours with respect to UEs density may be explained as follows:

- The first one corresponds to the case where UEs density is very low ($U=25$) or very high ($U=150$). In such case MLB impact is not tangible for the same reasons cited for Figure 9.

- The second one concerns the case where $U \in \{50, 75\}$. For such case although the MLB algorithm is very likely to be activated, the PLR remains almost unchanged. This may be explained by the relatively low UEs density not enough high to provoke packet loss.

- The third case corresponds to an UEs density around 100. In this particular case, MLB algorithm activation decreases the PLR value with respect to UEs density. This may be explained by both the relatively high UEs density and the highly probable MLB algorithm enabling.

Figure 11. PLR vs UEs density
Figure 12 illustrates the evolution of the number of successful HOs according to UEs density with and without MLB algorithm activation. The number of successful HOs is almost unchanged when activating MLB algorithm. This may be considered as a significant gain, since a MLB scheme promotes enabling more HOs in order to attenuate load disparity between cells. In conclusion the implemented MLB algorithm seems to be able to find convenient trade-off between different investigated network key performance indicators (KPIs) since it improves global network throughput and PLR without increasing the rate of HO signalling.

8. CONCLUSION

In this paper we propose an implementation, using ns-3, of a MLB algorithm based adaptive HO for downlink LTE SON systems. Since ns-3 with its current version (i.e. ns-3.24) did not implement yet some MM functions, we propose also a basic implementation of Resource Status Reporting EPs which is required for gathering information about network cell’s load conditions. Numerical results show that enhancements provided by the proposed MLB algorithm are closely dependent on network UEs density and it's illustrated in this paper through different KPIs. These enhancements concerns particularly PLR reduction, Jain’s fairness index and network global throughput increase without HO Overhead. In our future works we intend to further investigate and refine the implemented MM functions. More effort should be made to test and validate the proposed X2AP EPs extensions and to optimize MLB algorithm thresholds.

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